### Exercise 1.

Let  $f,g \in K[x]$ , the polynomial ring in *one* variable. Express the greatest common divisor and the least common multiple of f and g in terms of elements in syz(f,g) and derive an algorithm to compute these, assuming we can compute a Gröbner basis of syz(f,g).

**Solution:** Let d = gcd(f, g), then there exist  $k, l \in K[x]$  such that f = kd, g = ld. The element  $\binom{-l}{k}$  is then an element in syz(f, g) which determines the gcd of f and g uniquely (up to units).

Let c = lcm(f,g) then there exist  $a, b \in K[x]$  such that af = c, bg = c and the element  $\begin{pmatrix} a \\ -b \end{pmatrix}$  is in syz(f,g) and determines the lcm uniquely (up to units).

We show first, that for any  $f,g \in K[x]$  the syzygy module syz(f,g) is a rank one module, i.e., generated by one element.

Let  $\begin{pmatrix} \alpha \\ \beta \end{pmatrix} \in \operatorname{syz}(f,g)$ . Then

$$\begin{split} \alpha f &= -\beta g \\ \Rightarrow \alpha \frac{f}{\gcd(f,g)} &= -\beta \frac{g}{\gcd(f,g)} \end{split}$$

Since  $\frac{f}{\gcd(f,g)}$  and  $\frac{g}{\gcd(f,g)}$  are coprime, we see that  $\alpha$  is a multiple of  $\frac{g}{\gcd(f,g)}$  and  $\beta$  is a multiple of  $\frac{f}{\gcd(f,g)}$ . Since K[x] is a factorial ring, we can use primefactorization. Together with  $\alpha f = -\beta g$  it follows that the factor for both  $\alpha$  and  $\beta$  is the same.

$$\alpha = d \cdot \frac{g}{\gcd(f,g)}, \ \beta = -d \cdot \frac{f}{\gcd(f,g)}.$$

As  $\begin{pmatrix} \alpha \\ \beta \end{pmatrix}$  was a random element in syz(f,g) it follows that syz $(f,g) = \langle \begin{pmatrix} g \\ -\frac{g}{\gcd(f,g)} \\ -\frac{f}{\gcd(f,g)} \end{pmatrix} \rangle$ .

Furthermore, this generating set is a reduced Gröbner basis of syz(f,g) (since it consists of only one element). Thus, it is unique up to units.

Algorithm: INPUT:  $f, g \in K[\underline{x}]$ OUPUT: (Icm(f, g), gcd(f, g))INSTRUCTIONS:

• Compute a reduced Gröbner basis  $S = ( \begin{pmatrix} a \\ b \end{pmatrix} )$  of syz(f,g).

• 
$$\operatorname{lcm}(f,g) = af$$
,  $\operatorname{gcd}(f,g) = \frac{fg}{\operatorname{lcm}(f,g)}$ 

Return  $(\operatorname{lcm}(f,g), \operatorname{gcd}(f,g)).$ 

### Exercise 2.

Let

$$V = V((x+y) \cdot (x-y) \cdot (x+z^2)) \subset \mathbb{C}^3$$

and

$$W = V((x+z^2) \cdot (x-y) \cdot (z+y)) \subset \mathbb{C}^3.$$

Draw sketches of V and W in  $\mathbb{R}^3$ , and try to guess the ideal of  $\overline{V \setminus W}$ . How could you verify your guess? Answer with an approviate algorithm and describe the steps accurately without doing all the computations.

**Proof:** We rewrite V and W:





(a)  $V \cap \mathbb{R}^3$ 



Abbildung 1

$$V = V((x + y) \cdot (x - y) \cdot (x + z^2))$$
  
=  $V(x + y) \cup V(x - y) \cup V(x + z^2)$   
 $W = V((x + z^2) \cdot (x - y) \cdot (z + y))$   
=  $V(x + z^2) \cup V(x - y) \cup V(z + x)$ 

For the sketch see Figure 1.

The variety V(x + y) does not seem to be contained in W so we guess  $\overline{V \setminus W} = V(x + y)$ . We need to verify this claim by checking  $I = \sqrt{I}$  (1.5.13) and computing the ideal quotient (I : J) of

$$I = \langle (x+y) \cdot (x-y) \cdot (x+z^2) \rangle = \langle x^3 + x^2 z^2 - x y^2 - y^2 z^2 \rangle$$

and

$$J = \langle (x+z^2) \cdot (x-y) \cdot (z+y) \rangle = \langle x^2 z + 2y - xyz - xy^2 + z^3 + yz^2 - yz^3 - y^2 z^2 \rangle$$

Since I is a principal ideal, we know that for some  $r \in \sqrt{I}$  here is an  $n \in \mathbb{N}$ ,  $f \in K[x, y, z]$  such that  $r^n = f(x+y)(x-y)(x+z^2)$ . Since they are prime elements (x+y), (x-y),  $(x+z^2)$  all have to divide r, so  $r \in I$  and  $I = \sqrt{I}$ .

Computing the ideal quotient (I : J) using Variant 1 of algorithm 2.5.9, we first compute a Gröbner basis for  $I \cap J$  with algorithm 2.5.7. For that we have to eliminate t from

$$G = \langle t \cdot (x^3 + x^2 z^2 - xy^2 - y^2 z^2), (1 - t)(x^2 z + x^2 y - xyz - xy^2 + z^3 + yz^2 - yz^3 - y^2 z^2) \rangle.$$

We choose  $>_{lp}$  with t > x > y as our global elimination order and compute a Gröbner basis B of G. We write  $B = (g_1, ..., g_k)$  for some k and  $g = x^2z + y - xyz - xy^2 + z^3 + yz^2 - yz^3 - y^2z^2$  (then  $J = \langle g \rangle$ .) We then set  $g'_i = \frac{g_i}{g}$  and  $S = (g'_1, ..., g'_k)$ . This is a generating set of I : J by algorithm 2.5.9.

**Exercise 3.** a) Is the polynomial  $f = x^7 + x^2$  in the radical of  $I = \langle x^{11} + x^6, yx^4 + x^8, y^3 + x^2 \rangle$ ? Use algorithm 2.5.5 to prove your claim.

Hint: Eliminate y first.

**Solution:** Since  $f \in K[x]$ , we know that  $f \in \sqrt{I}$  if and only if  $f \in \sqrt{I \cap K[x]}$ .

We choose the global elimination monomial ordering  $>_{lp}$  with y > x and compute a grobner basis for I.

$$\begin{split} f_1 = x^{11} + x^0 \\ f_2 = yx^4 + x^8 \\ f_3 = y^3 + x^2 \\ P = \{(f_1, f_2), (f_1, f_3), (f_2, f_3)\} \\ \\ \text{Using product criterion we eliminate } (f_1, f_3) \\ & \text{spoly}(f_1, f_2) = y(x^{11} + x^6) - x^7(yx^4 + x^8) \\ & = -x^{15} + yx^6 \\ \\ \text{LM}(f_2)|\text{LM}(-x^{15} + yx^6) \\ q_2 = x^2 \\ r = yx^6 - x^{15} - x^2(yx^4 + x^8) \\ & = -x^{15} - x^{10} = -x^4(x^{11} + x^6) \\ \Rightarrow P = \{(f_2, f_3)\} \\ \text{spoly}(f_2, f_3) = y^2(yx^4 + x^8) - x^4(y^3 + x^2) \\ & = y^2x^8 - x^6 \\ \\ \text{LM}(f_2)|\text{LM}(y^2x^8 - x^6) \\ q_2 = yx^4 \\ r = y^2x^8 - x^6 - yx^4(yx^4 + x^8) \\ & = -x^6 - yx^{12} \\ q_2 = yx^4 - x^8 \\ r = -x^6 - yx^{12} + x^8(yx^4 + x^8) \\ & = -x^{16} + x^6 \\ q_1 = -x^5 \\ r = -x^{16} + x^6 + x^5(x^{11} + x^6) \\ & = x^6 + x^{11} = f_1 \\ \Rightarrow P = \emptyset \end{split}$$

So I is already given by a Gröbner basis. Hence  $I \cap K[x] = \langle x^{11} + x^6 \rangle$ . Now we apply Algorithm 2.5.5 to  $f = x^7 + x^2$  and  $J = \langle x^{11} + x^6 \rangle$ . We choose  $>_{lp}$  with t > x as the global elimination order. We need to compute a minimal Gröbner basis for

$$\begin{aligned} \langle x^{11} + x^6, 1 - t(x^7 + x^2) \rangle, & f_1 = x^{11} + x^6 \\ f_2 = 1 - t(x^7 + x^2) \\ S = \{f_1, f_2\} \\ P = \{(f_1, f_2)\} \\ \text{spoly}(f_1, f_2) = -t(x^{11} + x^6) - x^4(1 - t(x^7 + x^2)) \\ = -x^4 \\ \text{this cannot be divided by any Leading monomial in S so we add it to it.} \\ f_3 = -x^4 \\ S = \{f_1, f_2, f_3\} \\ P = \{(f_1, f_3), (f_2, f_3)\} \\ \text{spoly}(f_1, f_3) = -(x^{11} + x^6) + x^7(x^4) \\ = -x^6 \\ \text{LM}(f_3)|x^6 \Rightarrow q_3 = x^2 \\ r = -x^6 - x^2x^4 = 0 \\ \text{spoly}(f_2, f_3) = (-t(x^7 + x^2) + 1) + tx^3(x^4) \\ = -tx^2 + 1 \\ \text{this cannot be divided by any Leading monomial in S so we add it to it. \\ f_4 = -tx^2 + 1 \\ S = \{f_1, f_2, f_3, f_4\} \\ P = \{(f_1, f_4), (f_2, f_4), (f_3, f_4)\} \\ \text{spoly}(f_1, f_4) = t(x^{11} + x^6) + x^9(-tx^2 + 1) \\ = tx^6 + x^9 \\ \text{LM}(f_3)|tx^6 \\ r = tx^6 + x^9 - tx^2x^4 = x^9 = x^5(x^4). \\ \text{spoly}(f_3, f_4) = -(-tx^7 + x^2) + 1) + x^5(-tx^2 + 1) \\ = tx^2 - 1 + x^5 \\ \text{LM}(f_4)|tx^2 \\ r = tx^2 - 1 + x^5 - tx^2 + 1 = x^5 = x(x^4). \\ \text{spoly}(f_3, f_4) = -t(-x^4) + x^2(-tx^2 + 1) \\ = x^2 \\ \text{this cannot be divided by any Leading monomial in S so we add it to it. \\ f_5 = x^2 \\ S = \{f_1, f_2, f_3, f_4, f_5\} \\ P = \{(f_1, f_3), (f_2, f_3), (f_3, f_5), (f_4, f_5)\} \\ P = \{(f_1, f_5), (f_2, f_5), (f_3, f_5), (f_4, f_5)\} \\ P = \{(f_1, f_5), (f_2, f_2), (f_3, f_5), (f_4, f_5)\} \\ P = \{(f_1, f_2), (f_2, f_2) + 1 + tx^5(x^2) = -tx^2 + 1 = f_4 \\ \text{spoly}(f_3, f_3) = -tx^2 + 1 + t(x^2) = 1 \\ \Rightarrow \text{So we have to ad 1 to S.} \\ \end{aligned}$$

 $P=\emptyset.$  (here we use that by the product criterion with 1 there is always polynomial division with remainder 0

 $y^2$ )

### Introduction to Commutative Algebra and Algebraic Geometry Solution to Exercise Sheet 8

So  $S = \{f_1, f_2, f_3, f_4, f_5, 1\}$  is a Gröbner basis for  $\langle x^{11} + x^6, 1 - t(x^7 + x^2) \rangle$ . For the algorithm we need a minimal Gröbner basis, but we can already see that 1 will still be contained in the reduced Gröbner basis, so the output of the algorithm will be 1, i.e.,  $f \in \sqrt{I}$ .

b) Let 
$$I = \langle x^2 + 2y^2 - 3, x^2 + xy + y^2 - 3 \rangle \subset \mathbb{Q}[x, y]$$
. Compute  $I \cap \mathbb{Q}[y]$ .

**Solution:** We choose the gloabl elimination ordering  $>_{lp}$  with x > y. We compute a Gröbner basis of I.

$$\begin{split} f_1 &= x^2 + 2y^2 - 3 \\ f_2 &= x^2 + xy + y^2 - 3 \\ S &= \{f_1, f_2\} \\ P &= \{(f_1, f_2)\} \\ \text{spoly}(f_1, f_2) &= f_1 - f_2 = x^2 + 2y^2 - 3 - (x^2 + xy + y^2 - 3) \\ &= y^2 - xy \end{split}$$

this cannot be divided by any Leading monomial in S so we add it to it.

$$f_{3} = -xy + y^{2}$$

$$S = \{f_{1}, f_{2}, f_{3}\}$$

$$P = \{(f_{1}, f_{3}), (f_{2}, f_{3})\}$$
spoly $(f_{1}, f_{3}) = -yf_{1} - xf_{3} = -yx^{2} - 2y^{3} + 3y - x(-xy + y^{2})$ 

$$= -2y^{3} + 3y - xy^{2}$$

$$q_{3} = y$$

$$r = -xy^{2} - 2y^{3} + 3y - y(-xy + y^{2})$$

$$= -2y^{3} + 3y - y(y^{2})$$

$$= -3y^{3} + 3y$$

this cannot be divided by any Leading monomial in  ${\cal S}$  so we add it to it.

$$f_{4} = -3y^{3} + 3y$$

$$S = \{f_{1}, f_{2}, f_{3}, f_{4}\}$$

$$P = \{(f_{2}, f_{3}), (f_{1}, f_{4}), (f_{2}, f_{4}), (f_{3}, f_{4})\}$$
use Product criterion to delete  $(f_{1}, f_{4}), (f_{2}, f_{4})$ 

$$P = \{(f_{2}, f_{3}), (f_{3}, f_{4})\}$$
spoly $(f_{2}, f_{3}) = -yf_{2} - xf_{3} = -y(x^{2} + xy + y^{2} - 3) - x(-xy + y^{2})$ 

$$= -2xy^{2} - y^{3} + 3y$$

$$q_{3} = 2y$$

$$r = -2xy^{2} - y^{3} + 3y - 2y(-xy + y^{2})$$

$$= -y^{3} + 3y - 2y^{3} = -3y^{3} + 3y = f_{4}$$
spoly $(f_{3}, f_{4}) = -3y^{2}f_{3} + xf_{4} = -3y^{2}(-xy + y^{2}) + x(-3y^{3} + 3y)$ 

$$= -3y^{4} + 3xy$$

$$q_{3} = -3$$

$$r = -2x^{4} + 2ry + 2(-ry + y^{2})$$

$$r = -3y^{2} + 3xy + 3(-xy + y^{2})$$
  
=  $-3y^{4} + 3y^{2} = y(-3y^{3} + 3y) = yf_{4}$   
 $\Rightarrow P = \emptyset.$ 

So a Gröbner basis of I is given by  $S = \{x^2 + 2y^2 - 3, x^2 + xy + y^2 - 3, -xy + y^2, -3y^3 + 3y, \}$ . The only element of this Gröbner basis that has leading monomial in  $\mathbb{Q}[y]$  is  $-3y^3 + 3y$ . So  $I \cap \mathbb{Q}[y] = \langle -3y^3 + 3y \rangle$ .  $\Box$